

Day - 11 : Sequences and Convergence

Touseef Haider

Study Mathematics In Lean(MIL) Section 3.6

```
1 import Mathlib.Data.Real.Basic
2
3 def ConvergesTo (s : ℕ → ℝ) (a : ℝ) :=
4   ∀ ε > 0, ∃ N, ∀ n ≥ N, |s n - a| < ε
5
6 -- (s : ℕ → ℝ) is the first input. it is a sequence of real numbers.
7 -- (a : ℝ) is a real number that the sequence converges to.
8 -- rest is just convergence definition.
9
10 example : (fun x y : ℝ ↦ (x + y) ^ 2) = fun x y : ℝ ↦ x ^ 2 + 2 * x * y + y ^ 2 := by
11 -- fun x y : ℝ ↦ ... is how Lean writes lambda functions.
12 ext
13 -- ext is a tactic that stands for ''extensionality''. For functions, it means that to
14 -- prove two functions are equal, you need to prove they give the same output for every
15 -- possible input.
16 -- ring is a powerful tactic that can automatically prove identities that hold in any ring
17 ring
```

```
1
2 example (a b : ℝ) : |a| = |a - b + b| := by
3   -- congr tactic short for congruence tries to prove equalities by breaking them down. Here
4   -- , it sees that we are trying to prove |X| = |Y|. It will change the goal to proving X =
5   -- Y, if it can justify that X=Y implies |X| = |Y|.
6   congr
7   ring
```

```
1 example {a : ℝ} (h : 1 < a) : a < a * a := by
2   convert (mul_lt_mul_right _).2 h
3   -- mul_lt_mul_right is a theorem that essentially says if x<y and z>0 then x*z < y*z. (.)
4   -- .2 is selecting a specific implication.
5   -- We are also telling Lean we want to prove our goal by showing it is equivalent to h,
6   -- after multiplying both sides by some positive number.
7   -- convert tactic that asks Lean to change the goal to something else, provided we can
8   -- prove the new thing implies the old thing. It changes the goal to 1 * a < a * a, and
9   -- leaves us with two side goals:
10  -- 1. Prove that 1 * a = a
11  -- 2. Prove that a is positive.
12  · rw [one_mul]
13    -- rw [one_mul] address the first goal, It rewrites the term 1 * a to just a, because
14    -- multiplying by 1 does not change the value.
15  exact lt_trans zero_lt_one h
16  -- This addresses the second goal. It uses the transitivity of the less than relation to
17  -- show that since 0 < 1 and 1 < a, then 0 < a.
```

Below is the theorem that states that a constant sequence converges to the constant value.

```
1 theorem convergesTo_const (a : ℝ) : ConvergesTo (fun _ ↦ a) a := by
2
3   intro ε εpos
4   -- We need to prove ∀ ε > 0, ... . The intro tactic introduces variables for the universal
5   -- quantifiers. It introduces ε(a real number) and εpos(a proof that ε > 0). Our new goal
6   -- will be ∃ N, ∀ n ≥ N, |a - a| < ε.
7   use 0
8   -- We need to prove ∃ N, ... . The use tactic provides a witness for the existential
9   -- quantifier. Here, we use 0 as the witness for N. Our new goal is ∀ n ≥ 0, |a - a| < ε.
```

```

7   intro n nge
8   -- We need to prove  $\forall n \geq 0$ , ... We introduce n (a natural number) and nge (a proof that
   --  $n \geq 0$ ). Our new goal is  $|a - a| < \epsilon$ .
9   rw [sub_self, abs_zero]
10  -- sub_self rewrites  $a - a$  as 0. The goal becomes  $|0| < \epsilon$ .
11  -- abs_zero rewrites  $|0|$  as 0. The goal becomes  $0 < \epsilon$ .
12
13  apply  $\epsilon$ pos
14  -- apply tactic tries to prove the goal by applying a theorem or lemma. Here, it applies  $\epsilon$ 
   -- pos, which is a proof that  $\epsilon > 0$ . The goal is now proved.

```

Below is the theorem that states that the sum of two convergent sequences converges to the sum of their limits.

```

1 theorem convergesTo_add {s t :  $\mathbb{N} \rightarrow \mathbb{R}$ } {a b :  $\mathbb{R}$ }
2   (cs : ConvergesTo s a) (ct : ConvergesTo t b) :
3   ConvergesTo (fun n  $\mapsto$  s n + t n) (a + b) := by
4   intro  $\epsilon$ pos
5   -- Introduce  $\epsilon$  and  $\epsilon$ pos, which is a proof that  $\epsilon > 0$ . The goal is now to find an N such
   -- that for all  $n \geq N$ ,  $|s n + t n - (a + b)| < \epsilon$ .
6   dsimp -- this tactic unfolds the definition of ConvergesTo, making the goal clearer.
7   have  $\epsilon$ 2pos :  $0 < \epsilon / 2$  := by linarith -- This is a key step in many convergence proofs. We
   -- want to use the fact that s and t converge. We will need them to be closer than  $\epsilon/2$  to
   -- their limits.
8   --  $\epsilon$ 2pos is a proof that  $\epsilon / 2 > 0$ . We use linarith to derive this from  $\epsilon$ pos.
9   rcases cs ( $\epsilon / 2$ )  $\epsilon$ 2pos with  $\langle N_s, h_s \rangle$ 
10  -- This is where we use the hypothesis cs. Since s converges, the definition ConvergesTo s
   -- a holds for any positive number including  $\epsilon / 2$  (using  $\epsilon$ 2pos).
11  -- So cs ( $\epsilon / 2$ )  $\epsilon$ 2pos gives us a natural number  $N_s$  and a proof  $h_s$  that for all  $n \geq N_s$ ,  $|s$ 
   --  $n - a| < \epsilon / 2$ .
12  -- rcases with  $\langle N_s, h_s \rangle$  means we are unpacking the existential and universal quantifier.
   -- It extract the N (which it names  $N_s$ ) and the proof  $\forall n \geq N_s, |s n - a| < \epsilon / 2$  (which
   -- it names  $h_s$ ).
13
14  rcases ct ( $\epsilon / 2$ )  $\epsilon$ 2pos with  $\langle N_t, h_t \rangle$ 
15  -- This is similar to the previous step.
16  use max  $N_s N_t$ 
17  -- Now we need to provide our N for the sum sequence. We choose the maximum of  $N_s$  and  $N_t$ .
   -- This ensures that both conditions for s and t are satisfied for n greater than or equal
   -- to this N.
18  intro n nge
19  -- We introduce n and nge that  $n \geq \max N_s N_t$ . The goal is now to show that  $|s n + t n - ($ 
   --  $a + b)| < \epsilon$ .
20  have hns :  $n \geq N_s$  := le_of_max_le_left nge
21  -- We need to show that  $n \geq N_s$ . Since  $n \geq \max N_s N_t$ , and the maximum is always greater
   -- than or equal to its left component, we can conclude  $n \geq N_s$ . le_of_max_le_left is the
   -- theorem that says this, and nge is our proof that  $n \geq \max N_s N_t$ .
22  have hnt :  $n \geq N_t$  := le_of_max_le_right nge
23  calc
24     $|s n + t n - (a + b)| = |(s n - a) + (t n - b)|$  := by
25      congr; ring -- We rearrange the terms inside the absolute value. congr simplifies it
   -- to proving the inside is equal and ring prove that.
26     $\leq |s n - a| + |t n - b|$  := abs_add _ _ -- We use the triangle inequality (abs_add)
   -- which states  $|x + y| \leq |x| + |y|$  for any real numbers x and y.
27     $< \epsilon / 2 + \epsilon / 2$  := add_lt_add (h_s n hns) (h_t n hnt) -- We know from  $h_s$  that if  $n \geq N_s$ 
   -- (which hns proves), then  $|s n - a| < \epsilon / 2$ . So  $h_s$  n hns is the proof of  $|s n - a| < \epsilon /$ 
   -- 2. Similarly,  $h_t$  n hnt gives us  $|t n - b| < \epsilon / 2$ .
28    -- add_lt_add is a theorem that states if  $x < y$  and  $z < w$ , then  $x + z < y + w$ . We use it
   -- to combine the two inequalities.
29     $= \epsilon$  := by ring -- Finally, we use the ring tactic to simplify  $\epsilon / 2 + \epsilon / 2$  to  $\epsilon$ . This
   -- completes the proof that the sum of the sequences converges to the sum of their limits.

```

We are proving that if s converges to a, then $(c * s)$ converges to $(c * a)$

```

1 theorem convergesTo_mul_const {s :  $\mathbb{N} \rightarrow \mathbb{R}$ } {a :  $\mathbb{R}$ } (c :  $\mathbb{R}$ ) (cs : ConvergesTo s a) :
2   ConvergesTo (fun n  $\mapsto$  c * s n) (c * a) := by
3   -- First, let's handle the easy case where c is 0.
4   by_cases h : c = 0
5   -- This is the block for the c=0 case.

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6   -- convert changes the goal to proving that two things are equal.
7   -- Here, it changes the goal to proving 'convergesTo_const 0' is the same as our
    original goal.
8   · convert convergesTo_const 0
9   -- The first goal from 'convert' is to show '(fun n ↦ c * s n) = (fun n ↦ 0)
10  · rw [h] -- replace c with 0
11    ring -- The ring tactic can solve equations like 0 * s n = 0
12  -- The second goal is to show c * a = 0.
13  rw [h] -- replace c with 0
14  ring -- ring solve 0 * a = 0
15
16  -- This is the blok for  $c \neq 0$  case. h is now hypothesis  $c \neq 0$ .
17  have acpos : 0 < |c| := abs_pos.mpr h
18  -- our goal is to show ConvergesTo (fun n ↦ c * s n) (c * a)
19
20  -- Start the  $\epsilon$  - N proof
21  intro  $\epsilon$ _pos
22
23  -- Our hypothesis cs gives us an N for any positive epsilon.
24  -- We will use it with  $\epsilon' = \epsilon / |c|$ 
25  -- This is a valid positive epsilon because  $\epsilon > 0$  and  $|c| > 0$ .
26  have  $\epsilon'$ _pos :  $\epsilon / |c| > 0$  := div_pos  $\epsilon$ _pos acpos
27
28  -- Get the N from our hypothesis cs.
29  let ⟨ N, hN ⟩ := cs ( $\epsilon / |c|$ )  $\epsilon'$ _pos
30
31  -- Now we prove this N as our witness for the final goal.
32
33  use N
34
35  -- We now have to prove that for any  $n \geq N$ , the inequality holds.
36  intro n hn
37
38  -- The goal is now  $|(fun n ↦ c * s n) n - c * a| < \epsilon$ 
39  -- Let's simplify the expression
40  rw [← mul_sub, abs_mul]
41
42  -- The goal is now  $|c| * |s n - a| < \epsilon$ 
43  -- We can use the inequality from our hypothesis 'hN'
44  -- hN says that for  $n \geq N$ , we have  $|s n - a| < \epsilon / |c|$ 
45  -- So we can replace  $|s n - a|$  with something bigger to prove our goal
46  calc
47    |c| * |s n - a| < |c| * ( $\epsilon / |c|$ ) := by
48      -- this step is true because  $|s n - a| < \epsilon / |c|$ 
49      apply mul_lt_mul_of_pos_left
50      · exact hN n hn
51      · exact acpos
52    _ =  $\epsilon$  := by
53      -- This step is just algebra
54      rw [mul_div_cancel₀ _ (ne_of_gt acpos)]

```